



# On TTEthernet for Integrated Fault-Tolerant Spacecraft Networks

**Andrew Loveless**

**NASA Johnson Space Center (JSC)**

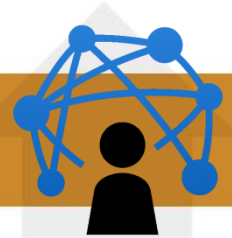
AIAA SPACE 2015, Aug. 31<sup>st</sup> – Sep. 2<sup>nd</sup> 2015

Pasadena, CA



# Project Overview and Motivation

- **Integrated modular avionics (IMA) principles are attractive for inclusion in spacecraft architectures.**
  - Consolidates multiple functions to shared computing platforms.
  - Reduces spacecraft cost, weight, and design complexity.
  - Interchangeable components increases overall system maintainability – important for long duration missions!
- **The Avionics and Software (A&S) project**
  - Funded by NASA's Advanced Exploration Systems program.
  - Developing a flexible mission agnostic spacecraft architecture according to IMA principles.
  - NASA can minimize development time and cost by utilizing existing commercial technologies.
  - Matures promising technologies for use in flight projects.



# Project Overview and Motivation

- **IMA Considerations in Networking**

- Requires network capable of accommodating traffic from multiple highly ***diverse systems*** (e.g. critical vs. non-critical) – potentially all from ***one shared computer platform***.
- Must prevent cascading faults b/w systems of differing criticalities connected to the same physical network.
  - ⚠ Most avionic system failures result from ineffective fault containment and the resulting domino effect.
- Some network technologies are better suited for certain tasks.
- Applying the same technology everywhere traditionally results in undue expense and limited performance.

Results in hybrid architectures with multiple technologies (e.g. NASA's LRO has MIL-STD-1553, SpaceWire, LVDS).

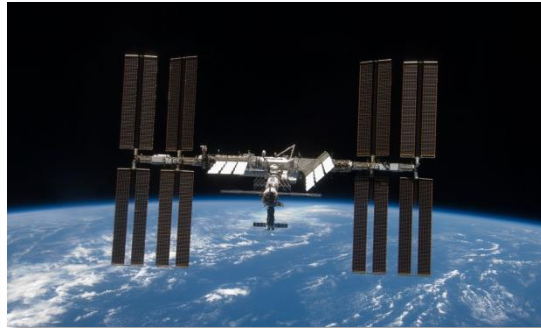


# Project Overview and Motivation

- **Ethernet is promising**
  - Inexpensive, widespread, and high speed = highly flexible.
  - Commonality promotes interchangeability between components.
  - Can augment with QoS enhancements for critical applications.
  - The A&S project considers Ethernet fundamental in the design of future manned spacecraft.
- **Integrated Power, Avionics, and Software (IPAS)**
  - Flexible evaluation environment for hardware and software in simulated mission scenarios.
  - Realistic framework of vehicle subsystems connected via Ethernet backbone.



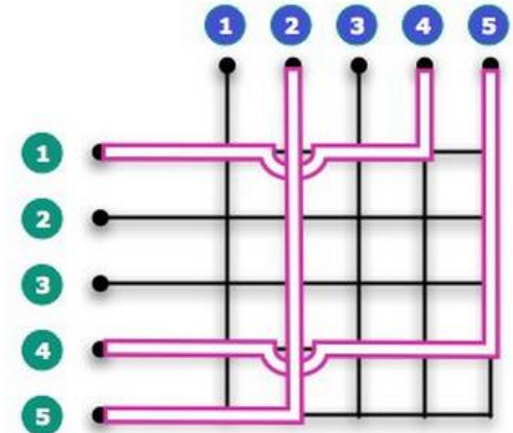
# Ethernet in Space Programs





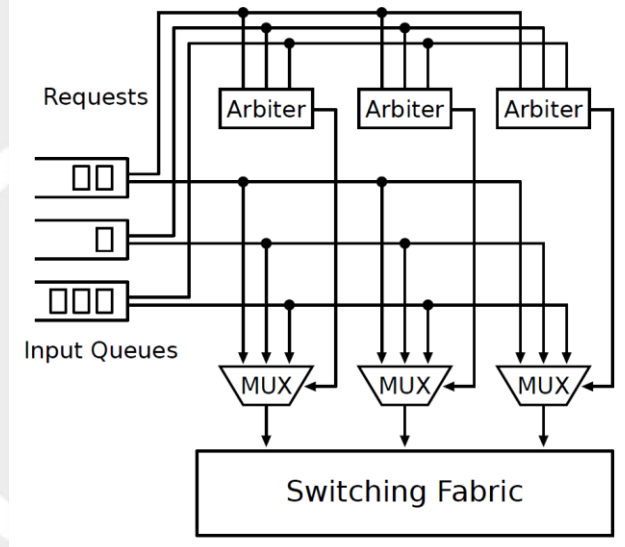
# Shortcomings of Classical Ethernet

- **Classical Ethernet characteristics**
  - Event-driven communication – messages are only sent in response to environmental or internal events (asynchronous).
  - Best-effort paradigm – no guarantees regarding transmission time or successful message delivery.
- **Timing within an Ethernet network is not predictable.**
  - Event-triggered = multiple frames will need to travel through the matrix simultaneously.
    - Usually supported by the switch fabric's parallel arrangement (space partitioning).
  - Collisions occur when frames are forwarded simultaneously to the same output port.
  - Arbitration is needed to regulate input to the switch fabric.



# Shortcomings of Classical Ethernet

- **What factors impact forwarding delay?**
  - 1) Degree of contention, 2) arbitration method
  - Frequency/severity of conflicts is highly variable.
- **Contention limits throughput**
  - Leads to buffer overflows and dropped frames.
  - 58.6% with input FIFOs under uniform traffic.
  - >80% with VOQs, crosspoint buffers, and better arbitration procedures (e.g. matrix, wavefront).
- **Modern advancements don't address unpredictable timing.**
  - E.g. VOQs eliminate head-of-line blocking, but still require arbitration.



Flight critical functions must operate in an entirely predictable manner and require a level of network determinism that classical Ethernet can't provide.

# Ethernet for Critical Applications

Quality of Service (QoS): Methods for controlling bandwidth, latency, jitter, or data loss in mission-critical networks (e.g. prioritization, traffic shaping).

- **“Industrial Ethernet” (e.g.  $\leq 100\text{Mbit/s}$  EtherNet/IP, PROFINET)**
  - Replaces proprietary Fieldbus solutions on factory floor (e.g. machinery).
  - Modified w/ master/slave arch., I/O controllers, and bus or ring topology.
  - RT services through specialized HW and extra protocols around payload.
- **Rate-Constrained (e.g. ARINC 664P7-1, IEEE 802.1BA AVB)**
  - Predetermined knowledge of traffic patterns (max size, frequency) ensures upper bound on TX delays.
  - A priori agreement of network devices prevents buffer overflows in switch.
  - Latency 1-10ms,  $< 500\mu\text{s}$  jitter, arbitration.







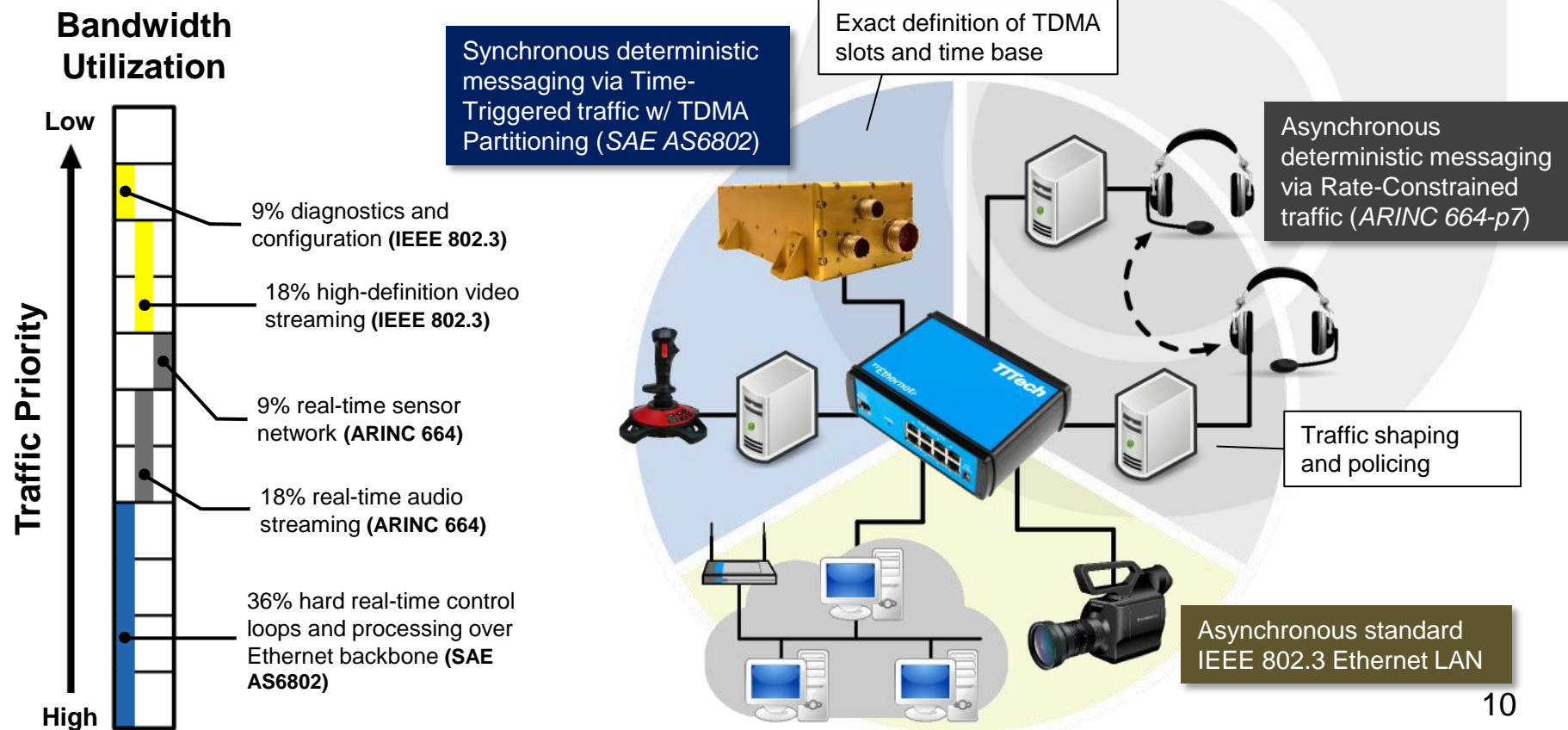
# Ethernet for Critical Applications

- **Time-Triggered Ethernet (SAE AS6802)**
  - Uses specialized end systems and network switches (like AFDX).
  - Network planning tool allocates each device a finite transmission window.
  - Each slot is repeated sequentially to form a periodic comm. schedule.
  - Config. files specifying schedule are loaded onto each network device.
- **Eliminating contention = no arbitration**
  - Decentralized synchronization process establishes a global time base.
  - Devices reference time to dispatch messages at predetermined instances.
  - Schedule guarantees no contention between TT frames.
  - Latency < 12.5  $\mu$ s/switch, < 1  $\mu$ s jitter, no arbitration

Note that controlling the jitter dramatically lowers latency compared to asynchronous RC traffic. A large portion of latency is the jitter!

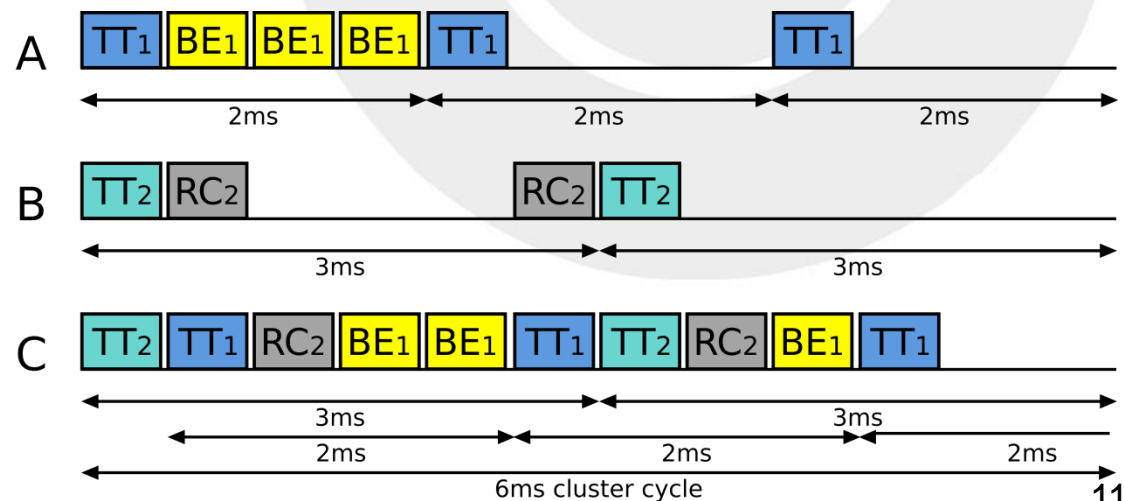
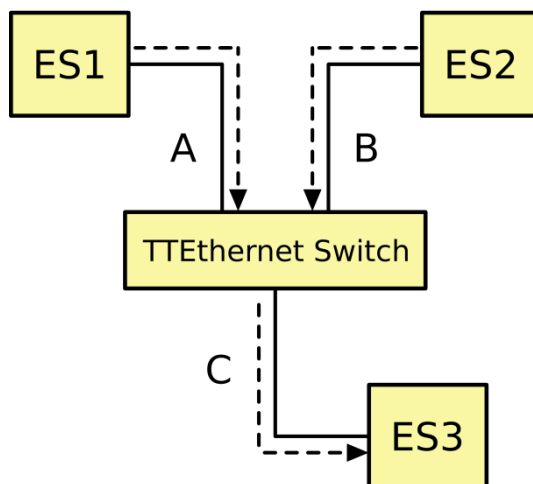
# TTEthernet Traffic Integration

TTEthernet overcomes difficulties in realizing an IMA architecture by providing three distinct traffic classes covering the full spectrum of criticality levels.



# TTEthernet Traffic Integration

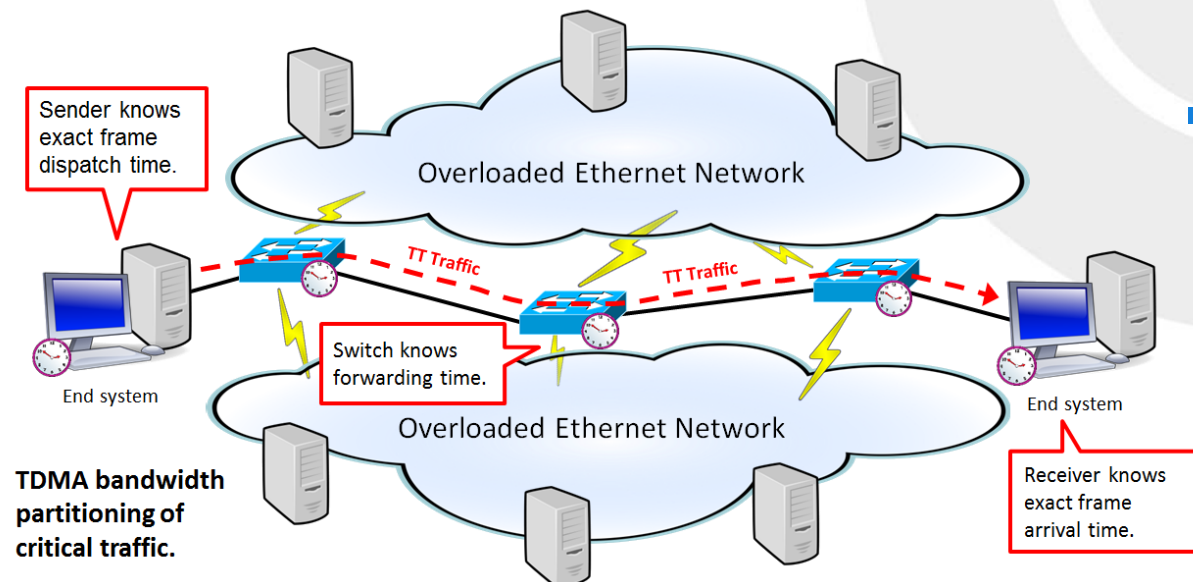
- **Priority-based partitioning: 3 traffic classes on 1 physical layer.**
  - Messages forwarded: 1) as scheduled (TT), or 2) as priority allows (RC, BE).
  - Bandwidth is released if TT message is not sent in synchronous time slot.
  - Ensuring determinism in a mixed-criticality network:
    - Timely block: Prevents RC or BE transmission during TT slots (unless freed).
    - Shuffling: Higher priority message is queued until lower priority frame is sent.



# TTEthernet Traffic Integration

TTEthernet network partitioning reduces cascading faults b/w platforms w/o the need for complex fault isolation procedures at the application level.

- **Traffic classes provide hard fault containment in the network.**
  - Guaranteed TT frame delivery regardless of asynchronous traffic patterns.
  - Communication schedule controls access of devices to network resources.



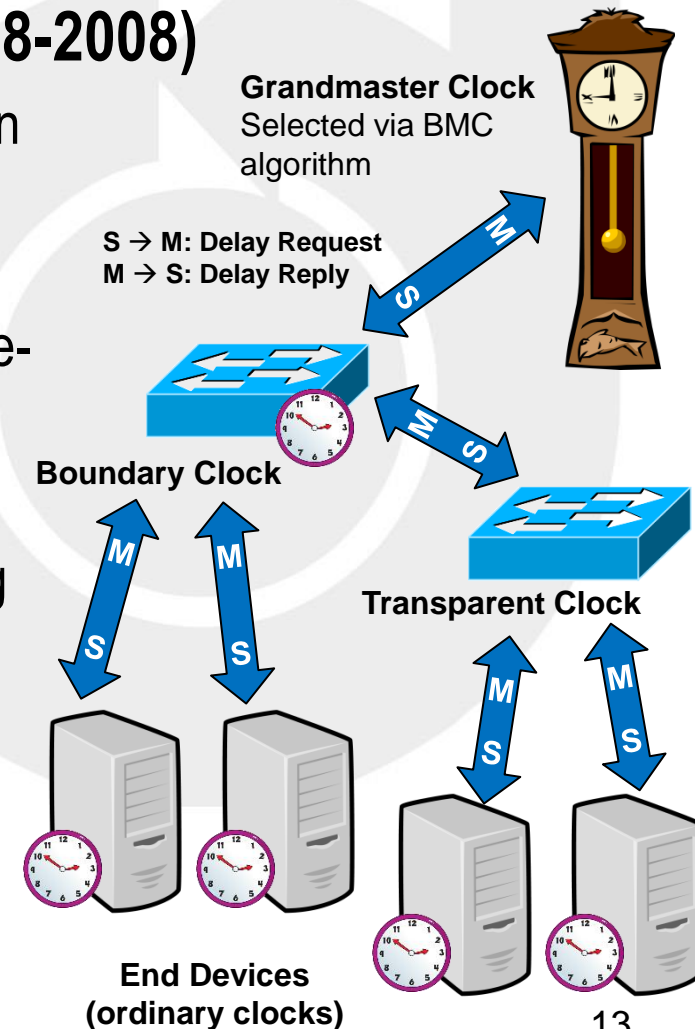
- Switches act as central bus guardians to protect against arbitrarily faulty end systems.

- TT: acceptance window
- RC: temporal distance

# Synchronization Comparison

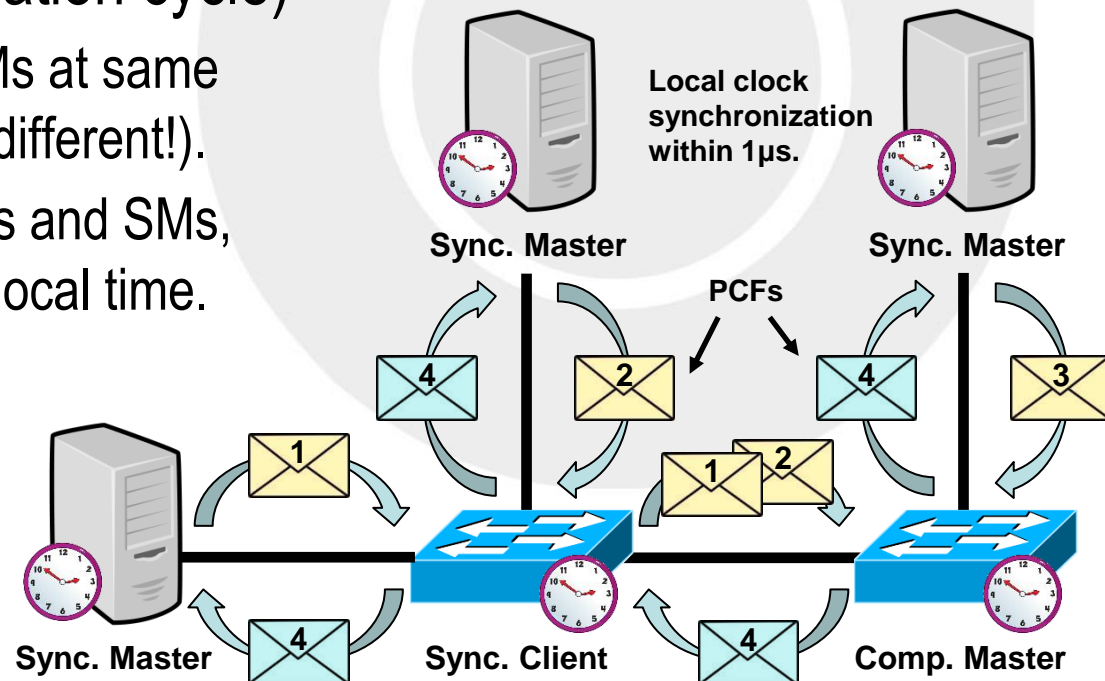
## • Precision Time Protocol (PTP IEEE 1588-2008)

- State-of-the-art Ethernet clock synchronization algorithm in industrial applications.
- Improves over Network Time Protocol (NTP) through specialized network hardware for time-stamping and decoding (sub- $\mu$ s accuracy).
- Protocol can be at Ethernet or IP layers.
- Hierarchical master/slave arch. for distributing time-of-day and clock frequency information.
- Uses best master clock (BMC) algorithm to select grandmaster clock source.
- Built-in redundancy means that if clock source fails, another is selected.



# Synchronization Comparison

- **Time-Triggered Ethernet (SAE AS6802)**
  - Based on the exchange of asynchronous Protocol Control Frames (PCFs).
  - Each component is assigned one of three roles (SC, SM, or CM).
- **Two Step Process (integration cycle)**
  - SMs dispatch PCFs to CMs at same local time (drift = actually different!).
  - CMs send PCFs to all SCs and SMs, which they use to correct local time.
- **Key Differences**
  - Decentralized “master”.
  - No search for best clock.
  - Tolerates multiple faults.
  - No external wall clock.





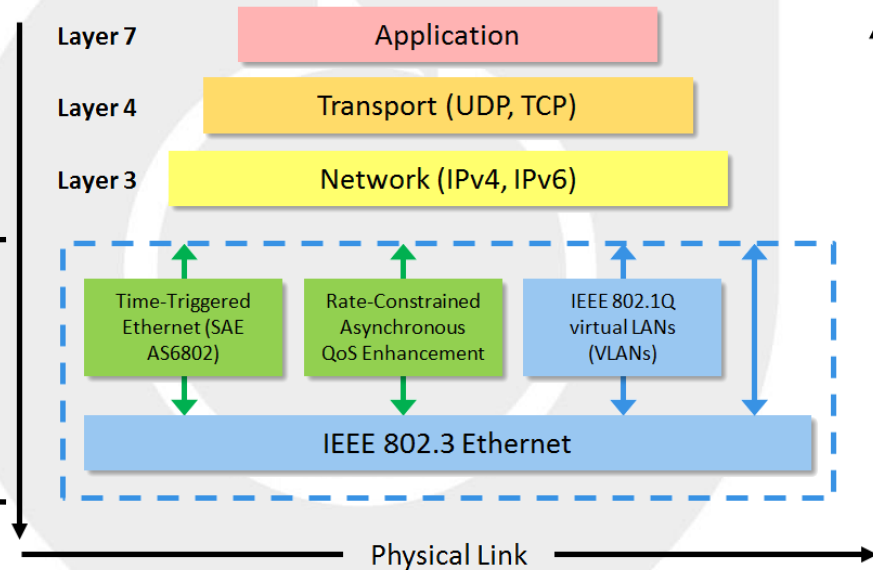
# TT/RC Network Stack Integration

- Directly alters Ethernet data link layer (L2). Does not add additional protocol layers.
- Traffic classes can coexist with other L2 QoS enhancements (e.g. IEEE 802.1Q).

Common higher level protocols (e.g. IPv4, UDP) can be used on top of TTEthernet's data link layer.

**TTEthernet**  
1Gbit/s Layer 2  
Ethernet Switch

## TCP/IP Model Network Stack



### IEEE 802.3 (Classical Ethernet)

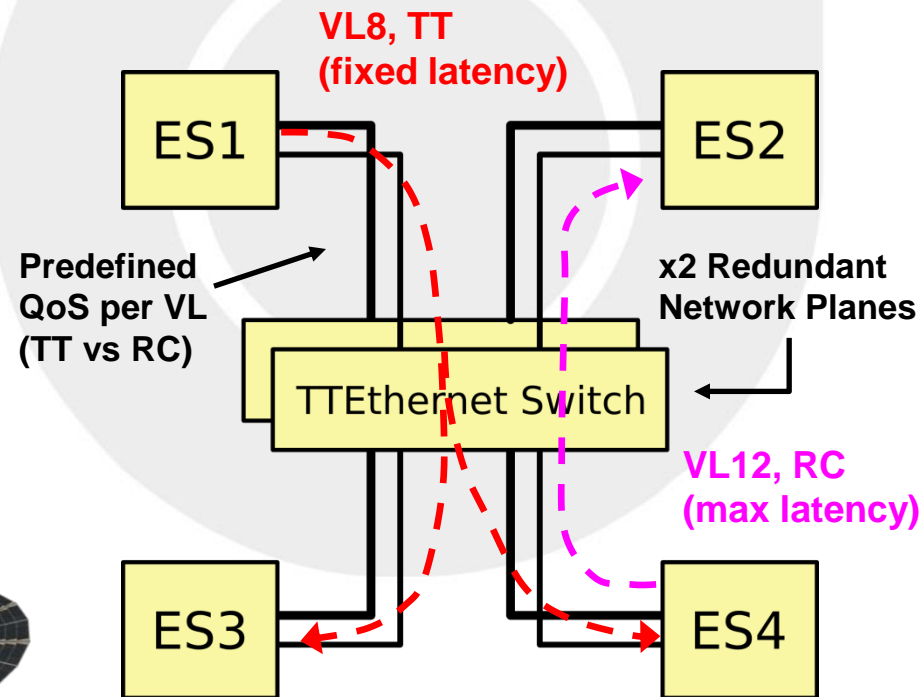
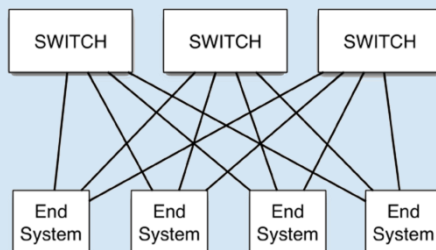
7 bytes	1 byte	6 bytes	6 bytes	2 bytes	46-1500 bytes	4 bytes	12 bytes
Preamble	SFD	Destination Address	Source Address	EtherType (Length)	Data Payload	FCS	IPG

### ARINC 664-P7 (RC) SAE AS6802 (TT)

7 bytes	1 byte	4 bytes	2 bytes	6 bytes	2 bytes	46-1500 bytes	4 bytes	12 bytes
Preamble	SFD	CT Marker	CTID	Source Address	Length	Data Payload	FCS	IPG

# Virtual Links and Redundancy

- **SAE AS6802 (TT) and ARINC 664p-7 (RC) use Virtual Links (VLs) to replace traditional MAC-based message delivery.**
  - Static forwarding table associates VLs with switch output ports.
  - VLs emulate point-to-point wiring seen in federated architectures.
- Increase fault-tolerance with multiple parallel switches.
- Redundancy mgmt. discards extra frames.
- Dual-fault tolerant w/ three redundant channels and high integrity devices.



Sample TTEthernet Network

# Flight Computer Failover

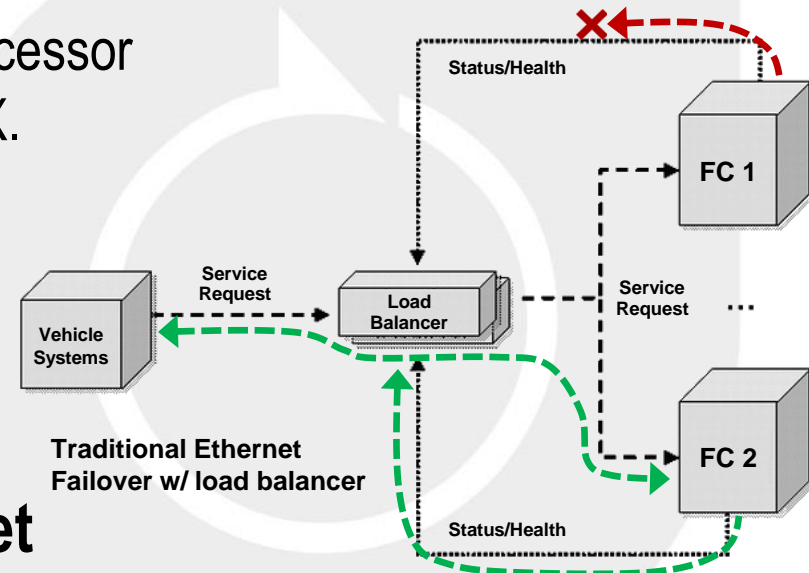
- **Past efforts used classical Ethernet over vehicle backbone.**

- Load balancer acted as virtual flight processor IP, detecting failure and directing TX/RX.
- Introduces single point of failure.
- Can increase fault tolerance w/ VRPP or redundant load balancers.

✗ Relies on monitoring with BE Ethernet.

- **Failover with deterministic Ethernet**

- Virtual link based delivery removes need for load balancer.
  - Identical messages can be dispatched to multiple recipients simultaneously.
- Means FC's have access to same data = More seamless failover.
- Can increase fault tolerance with redundant TTEthernet switches.
- Schedule driven communication compliments flight software behavior.



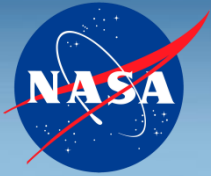


# Ascent Abort 2 (AA-2) Simulation


- **What is the Ascent Abort 2 Flight Test?**
  - Launch Abort System (LAS) carries CM away from ascent booster.
  - Goal is to stress the capabilities of synchronized redundant control loop.
  - Conducted AA-2 flight test demo in May '15 Integrated Test at JSC.
- **Redundant Flight Computer Architecture**
  - Three identical redundant flight computers (pc-linux).
  - Failover logic built into Core Flight Software System (CFS).
  - Synchronization over TTEthernet network (200Hz).
  - CFS included several genuine Orion fsw components:
    - Absolute Navigation (AbsNav) for Exploration Mission EM-1.
    - Service module abort, stochastic/optical navigation, and propellant balancing.
  - ANTARES simulation integrated into Tricksim.
    - Official NASA Orion spacecraft assessment tool used by JSC's GNC branch.



# Ascent Abort 2 (AA-2) Simulation



**trick**  **Simulation Environment**



LAS carries CM roughly 2 miles away from the launch vehicle at speeds up to 600 mph.

2

Attitude Control Motor (ACM) reorients CM to point heat shield forward/downward.

3

CM triggers abort event at altitude of maximum aerodynamic stress (Max Q). LAS separates CM from ascent booster.

1

LAS is separated from CM and jettisoned. LAS, CM, and booster free fall into the ocean.

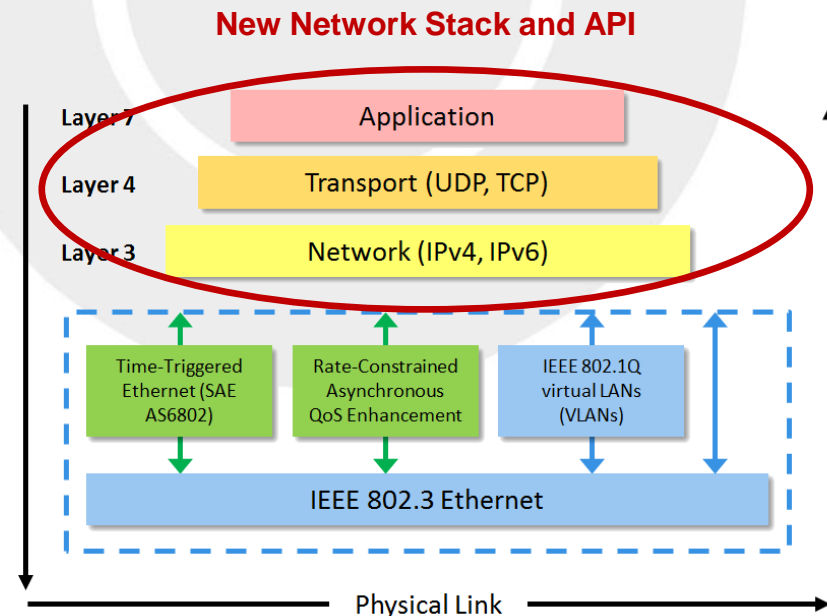
4

# Software-Level Network Stack

- **AA-2 – Unique Mission Requirements:**
  - Message payload sizes from simulation up to 20,000 bytes.
    - Ethernet frame data length is limited to 1500 bytes.
  - Throughput rates up to 100Mbit/s per Ethernet link.
  - Comm. with classical Ethernet systems w/o separate network adaptor.

- **Extension to TTEthernet Library (Phoenix IP - data link layer):**

- Implements IPv4 (RFC 791) and UDP (RFC 768) protocol layers.
- Abstraction from DMA management.
- Built in software = cross-platform.
- Maximizes throughput (e.g. minimize copies, parallel checksum summation).

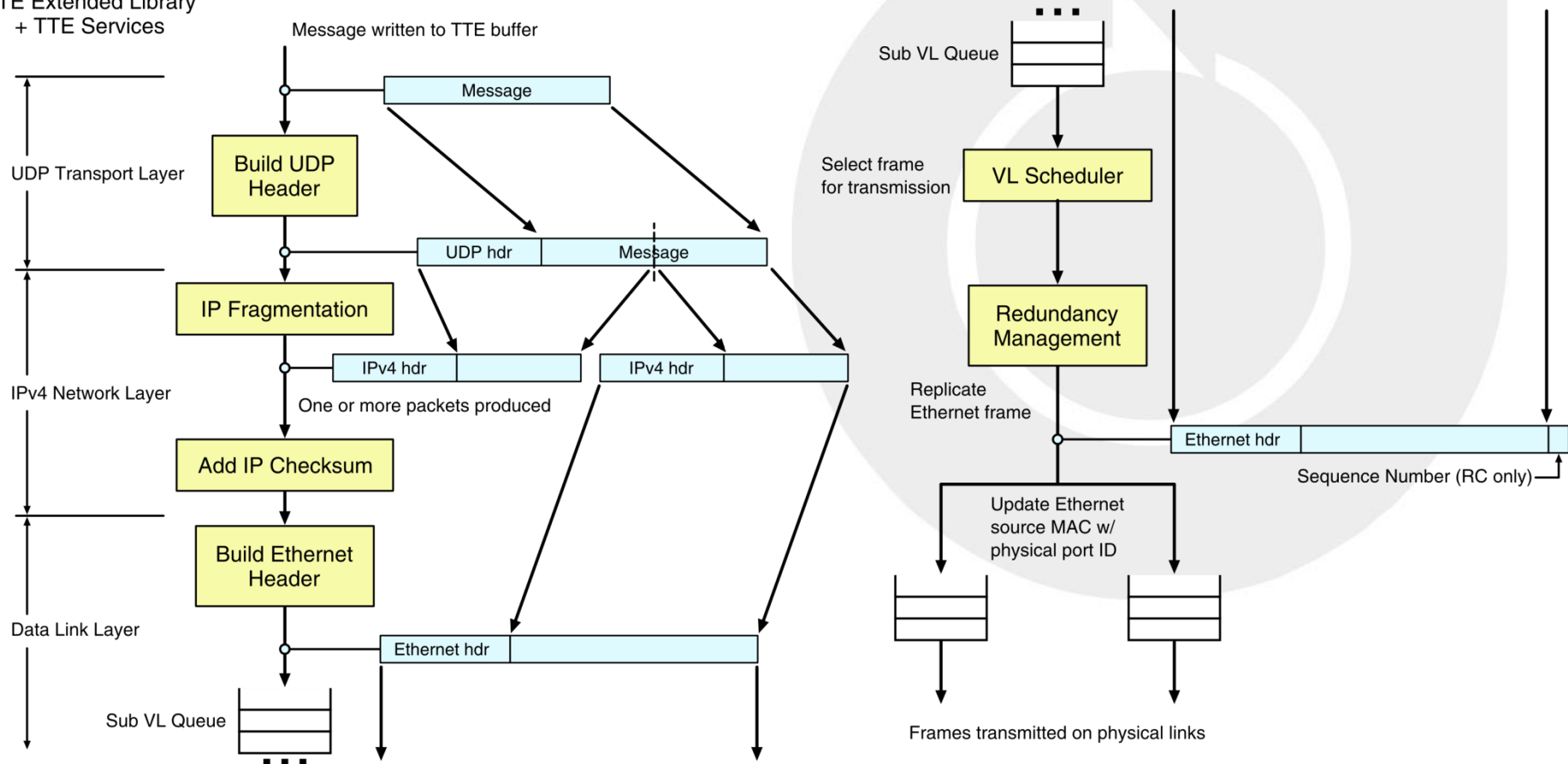




# Software-Level Network Stack

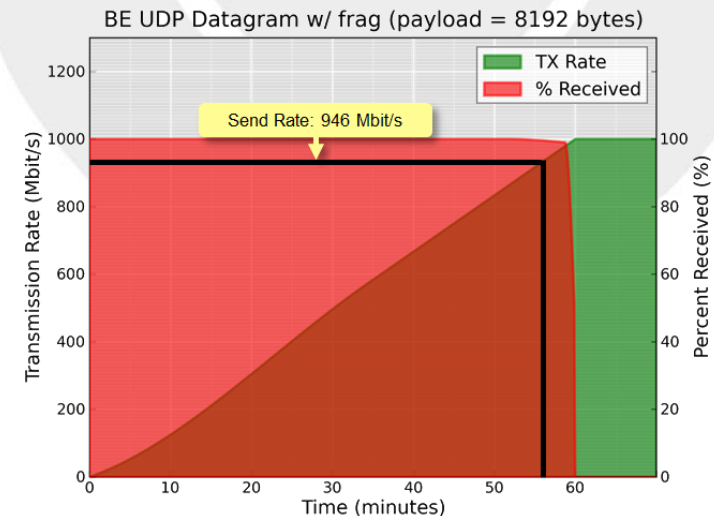
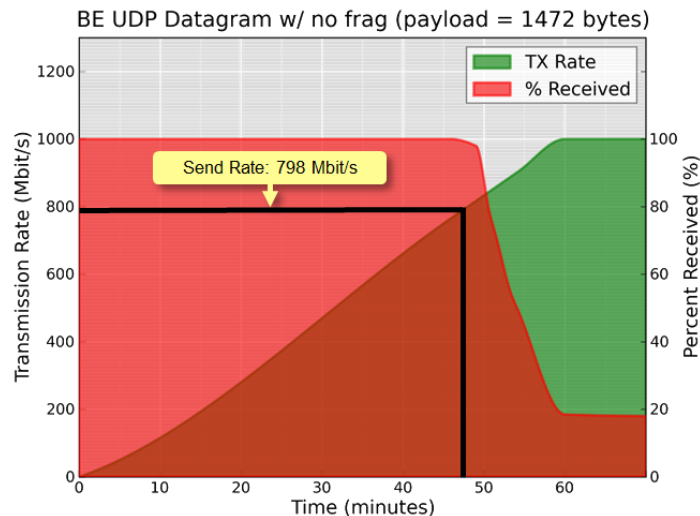
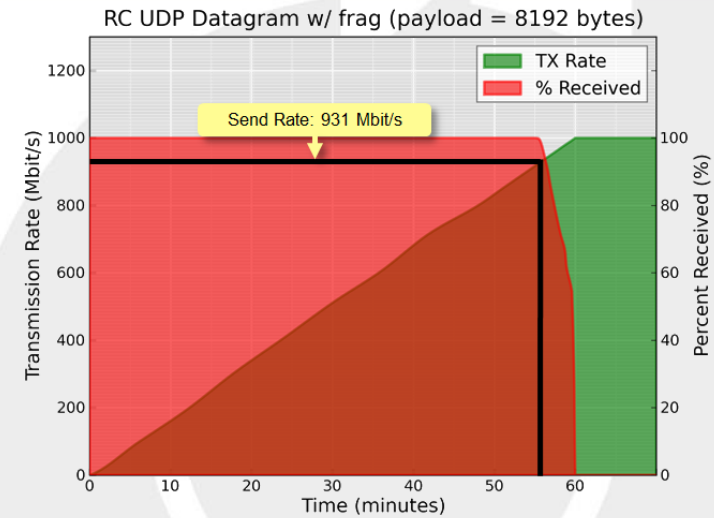
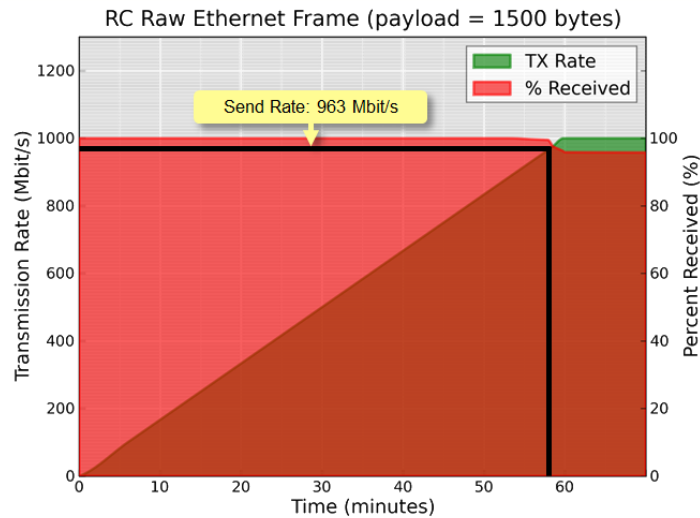
65,507 octets supported by library (max UDP data length according to RFC 5405)

TTE Extended Library  
+ TTE Services



TTEthernet Extended Library TX protocol stack

# Software-Level Network Stack

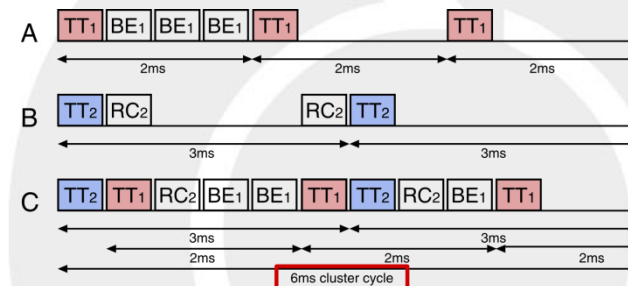
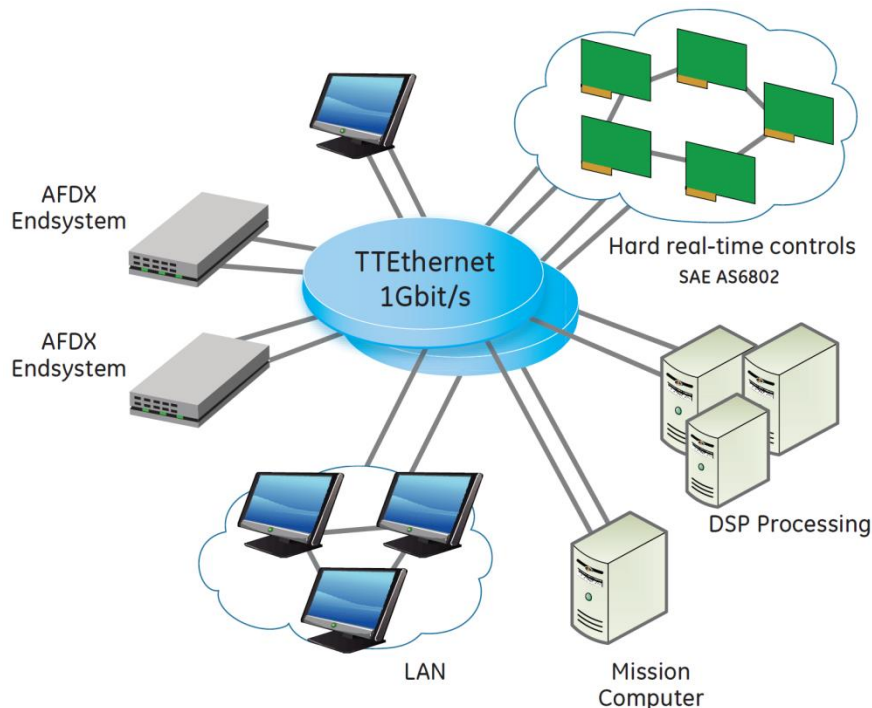


# Network-based CFS Scheduler



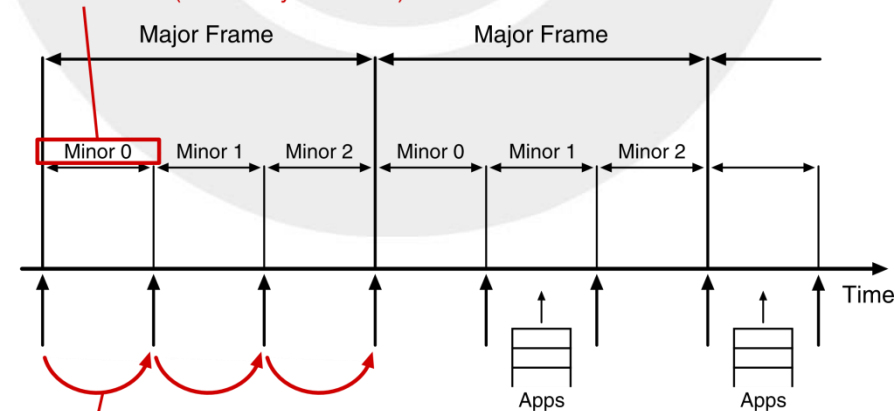
Combine the concept of scheduling the execution of CFS apps with the scheduling of the TTEthernet network.

- Drives FSX execution off cluster cycle.
- Can have deterministic scheduler even on limited hardware.



Cluster Cycle Period = LCM of all TT comm periods in sync domain

Minor Frame Period = (Cluster Cycle Period) x N



Trigger slot transition every N cluster cycle interrupts

# Flight Computer Synchronization

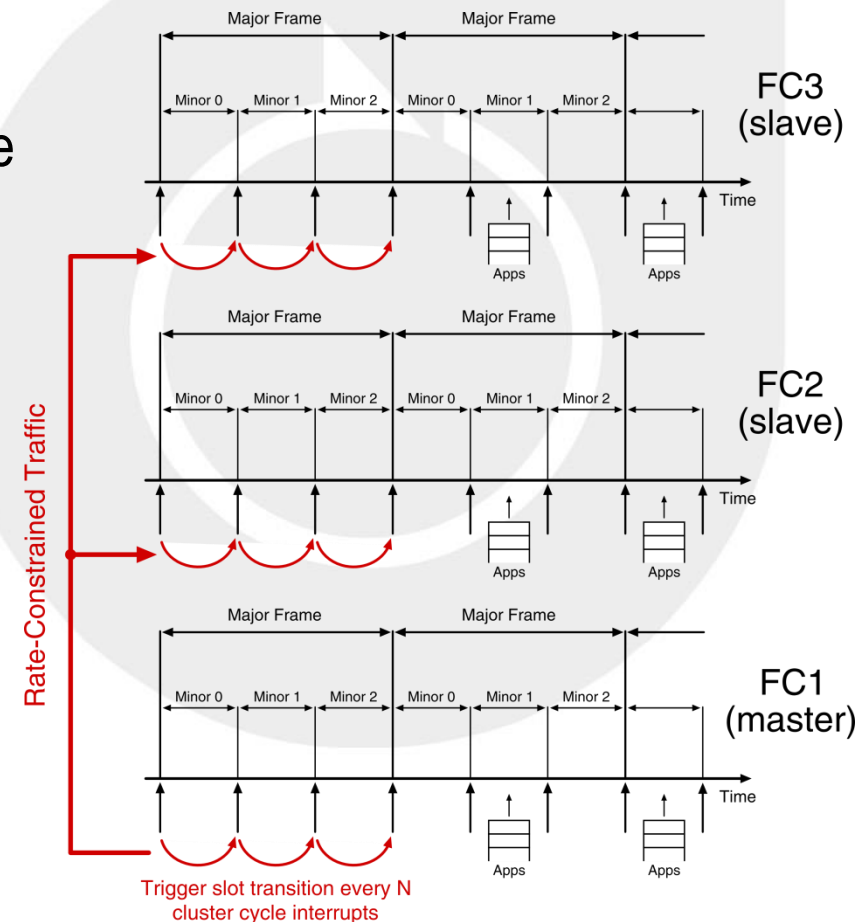


- **Message-based Synchronization**

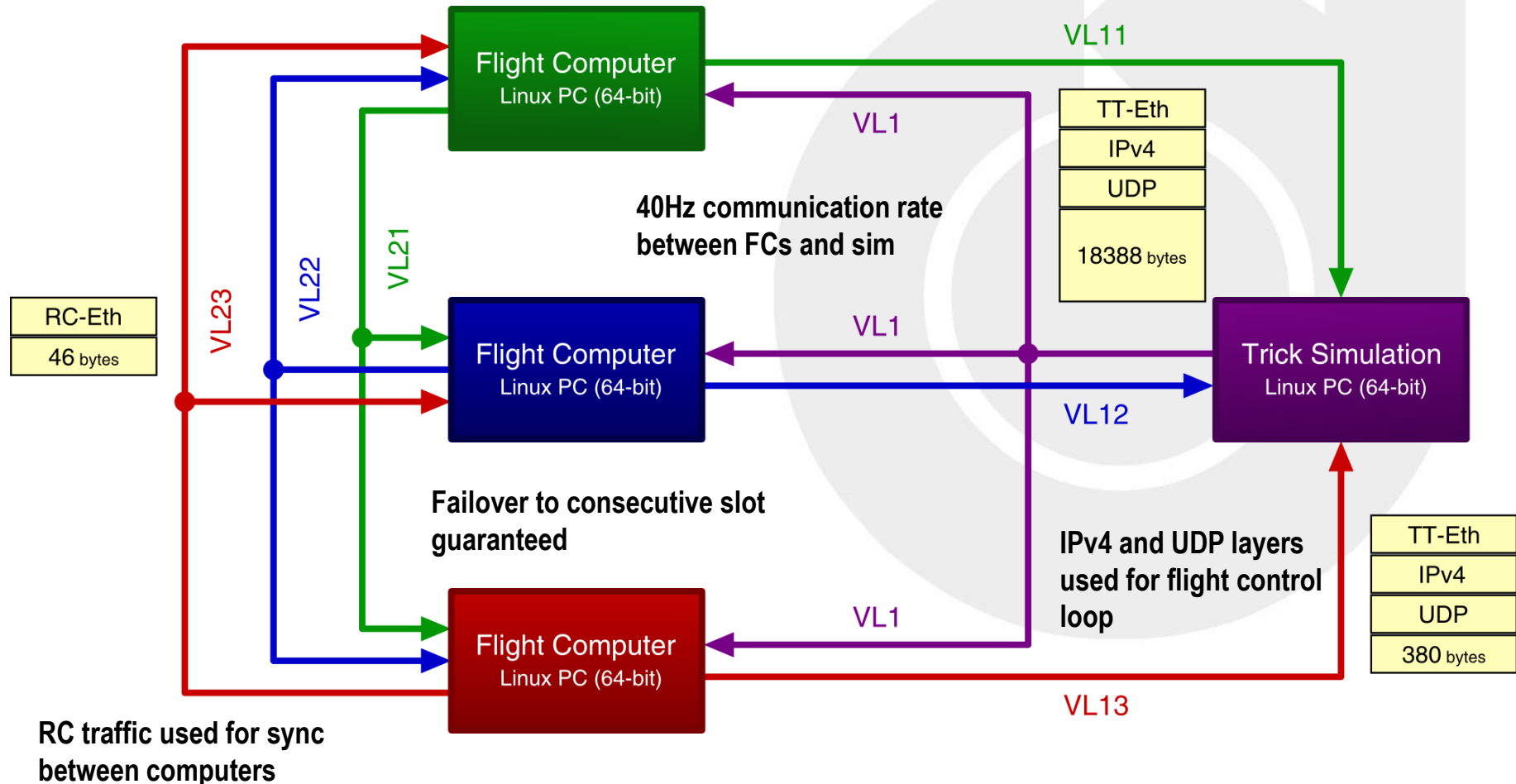
- Master/Slave architecture.
- Master computer drives CFS schedule off internal or network based timer.
- Highest-priority FC commands lower priority machines to move b/w slots.

- **Network-based Synchronization**

- Distributed architecture.
- Each FC drives CFS schedule off network interrupts (e.g. cluster cycle).
- Cluster period is a global property. Interrupts are generated on each machine simultaneously.



# Flight Computer Configuration





*Shaping the Future of Aerospace*